



中国科学院大学  
University of Chinese Academy of Sciences

CS101

# Systems Thinking

## Data Abstractions, Control Abstractions

zxu@ict.ac.cn  
zhangjialin@ict.ac.cn

# Outline

- What is systems thinking?
- Three objectives of systems thinking
- Abstraction
  - What is abstraction?
  - The COG properties of abstraction
  - One abstraction for many scenarios
  - Data abstractions
    - Representing numbers, characters
    - Review of bit, byte, character, integer, array, and slice, as well as struct
    - Pointers, files
  - Control abstractions
    - Precedence, sequence, selection, loop
    - Function and the four segments of text, data, stack, and heap
- Modularization
- Seamless transition

*These slides acknowledge sources for additional data not cited in the textbook*

## 3.4 Data abstractions

- Review data abstractions here in one place
  - To be more systematic
  - To understand why
- Representing numbers, characters
- Review of bit, byte, character, integer, array, slice
- Pointers, files

## 3.4.1 Positional number systems

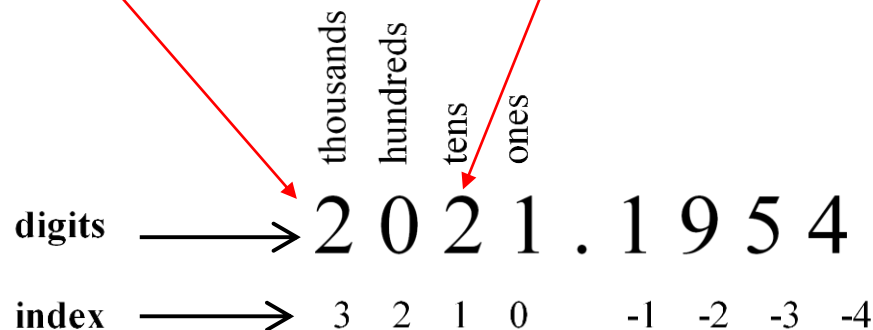
- Computers (and humans) use positional number systems
- Why?
  - What happens if we use Roman numerals?
    - a non positional number system, i.e., position independent system
- Q: MMXXI – MCMLIV = ?
  - Value of MMXXI = M + M + X + X + I = 1000+1000+10+10+1=2021
  - Value of MCMLIV = M + CM + L + IV = 1000+900+50+4=1954
    - Note that CM and IV are short-hand symbols, e.g, IV = IIII = I+I+I+I
- A: Use decimal, 2021-1954 = 67
- Do MMXXI – MCMLIV = LXVII without decimal arithmetic
- How about

- MMXXI + MCMLIV = ?
- MMXXI × MCMLIV = ?
- MMXXI ÷ MCMLIV = ?

Roman	M	D	C	L	X	V	I	IV	IX	XL	XC	CD	CM
Decimal	1000	500	100	50	10	5	1	4	9	40	90	400	900

# Positional number systems

- Positional number systems make arithmetic much easier
- Consider decimal number  $a = 2021.1954$
- Value  $a = \sum_{i=-4}^3 (a_i \times 10^i)$ , where
  - 10 is **base**,  $i$  is **index**,  $\{0,1,2,3,4,5,6,7,8,9\}$  is **digit set**
- The key point: The value of a digit depends on both the digit and the position (index) of the digit
  - The first 2 is at position 3, and the second 2 is at position 1
  - The first 2 represents  $a_3 \times 10^3 = 2 \times 10^3 = 2000$
  - The second 2 represents  $a_1 \times 10^1 = 2 \times 10^1 = 20$



# Examples of positional number systems

- Consider any n-digit number  $a = \sum_{i=0}^{n-1} (a_i \times b^i)$ , where
  - $b$  is the **base**,  $i$  is the **index**,  $\{0, \dots, b-1\}$  is the **digit set**
- There are three positional number systems often used
- Binary
  - $a = \sum_{i=0}^{n-1} (a_i \times 2^i)$ ,  $b = 2$ , digit set =  $\{0, 1\}$
  - Used in computers. This is what computers can understand
- Decimal
  - $a = \sum_{i=0}^{n-1} (a_i \times 10^i)$ ,  $b = 10$ , digit set =  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
  - Used by humans and high-level language programs
- Hexadecimal
  - $a = \sum_{i=0}^{n-1} (a_i \times 16^i)$ ,  $b = 16$ ,  
digit set =  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$
  - Used by humans and high-level language programs

# Are there other positional number systems?

- Radically different
  - Can we use an irrational number as the base, e.g.,  $\pi$ ?
- Yes
- Examples
  - Bergman's number system (the  $\tau$  number system)
    - Can represent some irrational numbers exactly in finite digits
    - Created by George Bergman, a 12-year junior high school student
    - Base  $b = \tau = (1 + \sqrt{5})/2 \approx 1.6180339$ , digit set =  $\{0, 1\}$
    - $14 = \mathbf{100100.110110} = \tau^5 + \tau^2 + \tau^{-1} + \tau^{-2} + \tau^{-4} + \tau^{-5}$
  - Fibonacci number system
    - Fibonacci numbers 1, 2, 3, 5, 8, ... as positional weights
    - digit set =  $\{0, 1\}$
    - $14 = \mathbf{11001} = 1 \times 8 + 1 \times 5 + 0 \times 3 + 0 \times 2 + 1 \times 1$

# Five positional number systems

Decimal	Hexadecimal	Binary	The $\tau$ Number System	FNS
$10^1 10^0$	$16^0$	$2^3 2^2 2^1 2^0$	$\tau^5 \tau^4 \tau^3 \tau^2 \tau^1 \tau^0 \tau^{-1} \tau^{-2} \tau^{-3} \tau^{-4} \tau^{-5} \tau^{-6}$	8 5 3 2 1
0	0	0000	0	00000
1	1	0001	1	00001
2	2	0010	10.01	00010
3	3	0011	100.01	00100
4	4	0100	101.01	00101
5	5	0101	1000.1001	01000
6	6	0110	1010.0001	01001
7	7	0111	10000.0001	01010
8	8	1000	10001.0001	10000
9	9	1001	10010.0101	10001
10	A	1010	10100.0101	10010
11	B	1011	10101.0101	10100
12	C	1100	100000.101001	10101
13	D	1101	100010.001001	11000
<b>14</b>	<b>E</b>	<b>1110</b>	<b>100100.110110</b>	<b>11001</b>
15	F	1111	100101.001001	11010



## 3.4.2 IEEE 754

### Use floating-point numbers to represent reals

- How to represent  $\pi \approx 3.1415927$  ?

- A simple way: converting whole and fraction into binary

- $3.1415927 = 11.0010010000111111011011$

- Scientific notation:  $3.1415927 =$

- Not unique:

- $3.1415927 = 31415927 \times 10^{-7} = 0.31415927 \times 10^1 = \dots$

$$= \underset{\substack{\uparrow \\ \text{Significant}}}{3.1415927} \times 10^{\underset{\substack{\downarrow \\ \text{Exponent}}}{0}}$$

- Key innovations of IEEE 754

- Normalized significant* to guarantee representation uniqueness

- Default left-most 1: assume one and only one 1-digit before the binary point

- Since it is default, the left-most 1-bit can be omitted, saving one bit

- Biased exponent* to speed up exponent comparison

- Special values* for unusual

- Infinities ( $\pm\infty$ ),
- Subnormal (underflow) values
- Not a Number (NaNs, such as trying to find  $\sqrt{-5}$ )

$$\begin{aligned} \pi &\approx 3.1415927 \times 10^0 \approx 1.5707964 \times 2^1 \\ &\approx +1.10010010000111111011011 \times 2^{00000001}; \\ &\rightarrow +.10010010000111111011011 \times 2^{00000001}; \quad \text{omit default left-most 1} \end{aligned}$$

Sign      Exponent      Significant

$$= \text{01000000010010010000111111011011}; \text{ the IEEE 754 representation}$$

# Floating-point numbers are approximate values

- How to test the equality of two floating-point numbers?

> go run ./testPoint123.go

0.1+0.2 == 0.3

0.1+0.2 != 0.3

0.1+0.2 == 0.3

>

- To test if (X+Y) is equal to Z
  - Don't use (X+Y)==Z
  - Test if the absolute value of the difference is less than epsilon, i.e., use  $\text{Abs}(X+Y-Z) < 10^{-12}$

There may be a small difference between 0.1+0.2 and 0.3

```
package main // testPoint123.go
import "fmt"
import "math"
func main() {
    if 0.1 + 0.2 == 0.3 {
        fmt.Println("0.1+0.2 == 0.3")
    } else {
        fmt.Println("0.1+0.2 != 0.3")
    }
    X := 0.1 // var X float64 = 0.1
    Y := 0.2
    Z := 0.3
    if X + Y == Z {
        fmt.Println("0.1+0.2 == 0.3")
    } else {
        fmt.Println("0.1+0.2 != 0.3")
    }
    if math.Abs(X+Y - Z) < math.Pow(10, -12) {
        fmt.Println("0.1+0.2 == 0.3")
    } else {
        fmt.Println("0.1+0.2 != 0.3")
    }
}
```

## 3.4.3 Review ASCII, Unicode, UTF-8

- ASCII encodes English characters, using 7 bits
- Unicode encodes the world's characters
  - using  $0000\sim FFFF_{16}$ ,  $10000\sim 10FFFF_{16}$ 
    - more than 1 million code points, some combinations reserved
    - A character's Unicode encoding needs at least 16 bits
- UTF-8 (Unicode Transformation Format – 8-bit)
  - A variable-width character encoding implementing Unicode
  - Historical standard; most widely used in Internet

Symbol	Description	ASCII	Unicode	UTF-8	Bytes needed by UTF-8
T	English capital letter T	0X54	U+0054	0X54	1
Ω	Greek letter Omega	N/A	U+03A9	0XCEA9	2
€	The Euro sign	N/A	U+20AC	0XE282AC	3
志	A Chinese character	N/A	U+5FD7	0XE5BF97	3
⓪	A Gothic letter	N/A	U+10348	0XF0908D88	4

# 3.4.4 Review of bit, byte, character, integer, array, slice

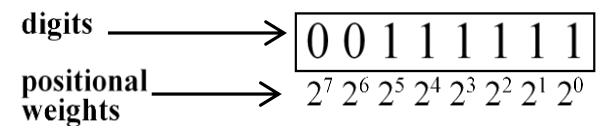
- Use a program to understand the basics in one place

```
X := byte(63)           // X is a byte variable. What if changed to X:=63?
fmt.Printf("Decimal: %d\n", X) // Decimal: 63
fmt.Printf("Hex: %X\n", X)   // Hex: 3F
fmt.Printf("Character: %c\n", X) // Character: ?
fmt.Printf("Binary: %b\n", X) // Binary: 111111
```

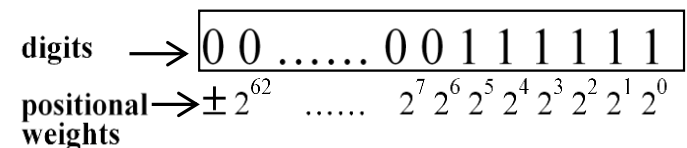
```
var S [5]byte = [5]byte{'h','e','l','l','o'}
           // S=[104, 101, 108, 108, 111]
var byteSlice []byte = S[1:4]
           // byteSlice=[101, 108, 108]
           // slice is built from array
```

```
fmt.Println("array S = ", S)
           // Array S = [104 101 108 108 111]
fmt.Println("byteSlice = ", byteSlice)
           // byteSlice = [101 108 108]
```

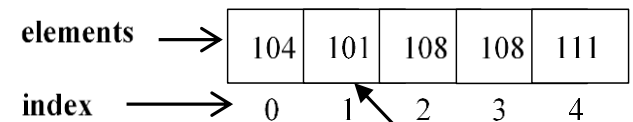
A **byte** variable X  
by X:=byte(63)



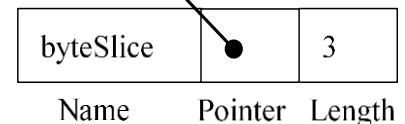
An **int** variable X  
generated by X:=63



An **array** S generated  
by var S [5]byte =  
[5]byte{'h','e','l','l','o'}



A **slice** byteSlice  
generated from array S  
var byteSlice []byte = S[1:4]



# The struct type

- Array is simple
  - Linear, consecutive arrangement of N elements of the same type
  - Example: `var A [5]byte` defines a byte array A of 5 elements
    - `A[0]`, `A[1]`, `A[2]`, `A[3]`, `A[4]` are all of type `byte`
- What if the elements are of different types?

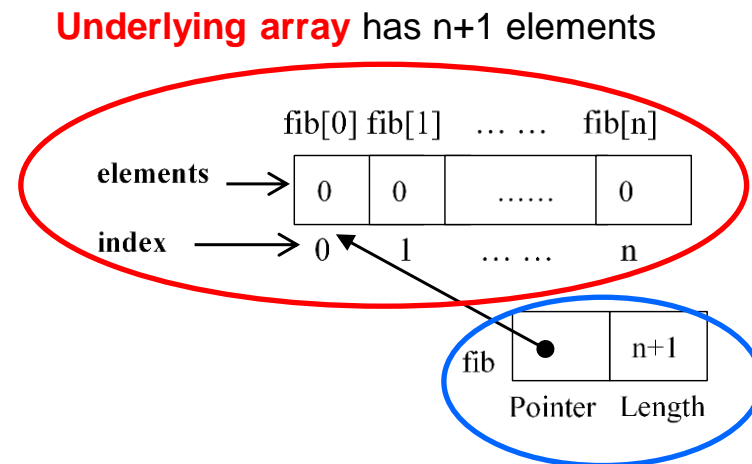
- Use struct

- A data structure with different types of elements
- Elements are called fields
- Use the dot notation to access fields, e.g.,
  - `JoanSmith.ID` accesses the student ID of Joan Smith
  - `FanWang.active = false` assigns the false value to the active field of student Fan Wang, indicating that he is not actively enrolled

```
type Student struct {  
    name      string  
    ID        int  
    majorCode byte  
    active    bool  
    contact   string  
}  
var JoanSmith, FanWang Student
```

# Remarks

- byte and int are the types most often used
  - Use int to represent integers
  - Use int to represent characters, assuming leading 0's
    - '?' in previous slide
  - Use byte (which is the same as uint8) to represent ASCII characters and small natural numbers
- A slice can be generated in two ways
  - From an existing array, e.g.,
    - `var byteSlice []byte = S[1:4]`
  - Using the make function
    - `fib := make([]int, n+1)`  
creates **a slice** and assigns to fib
    - `fib[i]` accesses the *i*th element of the array, which is initialized to 0. Also, `len(fib)=n+1`



## 3.4.6 How to do bit operations?

- Answer: Operate on the byte or int variable containing the bit
  - Through some mask mechanism
- Example 1: Inverting the rightmost bit of a byte
  - Input: 00111111<sup>1</sup>; Output: 00111111<sup>0</sup>
- Code explained with the corresponding operations

```
x := byte(63)      // assign 63=00111111 to variable x
v := ^x             // bitwise NOT of x, i.e., v=11000000
v = v & 0x1         // bitwise AND to retain the right-most bit of v
x = x & 0xFE        // bitwise AND to clear the right-most bit of x
x = x | v           // bitwise OR to get the final result
```

Mask  
mechanism

```
x = 001111111
v =  $\overline{0}\overline{0}\overline{1}\overline{1}\overline{1}\overline{1}\overline{1}\overline{1}$  = 110000000
v = 11000000 & 00000001 = 000000000
x = 00111111 & 11111110 = 001111100
x = 00111110 | 00000000 = 001111100
```

Given input  
Bitwise NOT  
Bitwise AND  
Bitwise AND  
Bitwise OR

# The code can be simplified

- Example 1: Inverting the rightmost bit of a byte
  - Input: 0011111**1**; Output: 0011111**0**

```
x := byte(63)      // assign 63=00111111 to variable x
x = x ^ 0x1         // bitwise XOR of x and 00000001
                    // i.e., 00111111 ^ 00000001 = 00111110
```

- However, the Masking mechanism is more general-purpose



# Replacing the least significant 2 bits of a byte

(used in Project Text Hider)

- Input: 001111**11**, 001010**10**; Output: 001111**10**
- Code explained with the corresponding operations

```
x := byte(63)      // assign 63=00111111 to variable x
v := byte(42)      // assign 42=00101010 to variable v
v = v & 0x3         // bitwise AND to retain the right-most 2 bits of v
x = x & 0xFC        // bitwise AND to clear the right-most 2 bits of x
                    // and retaining the leftmost 6 bits
x = x | v           // bitwise OR to get the final result
```

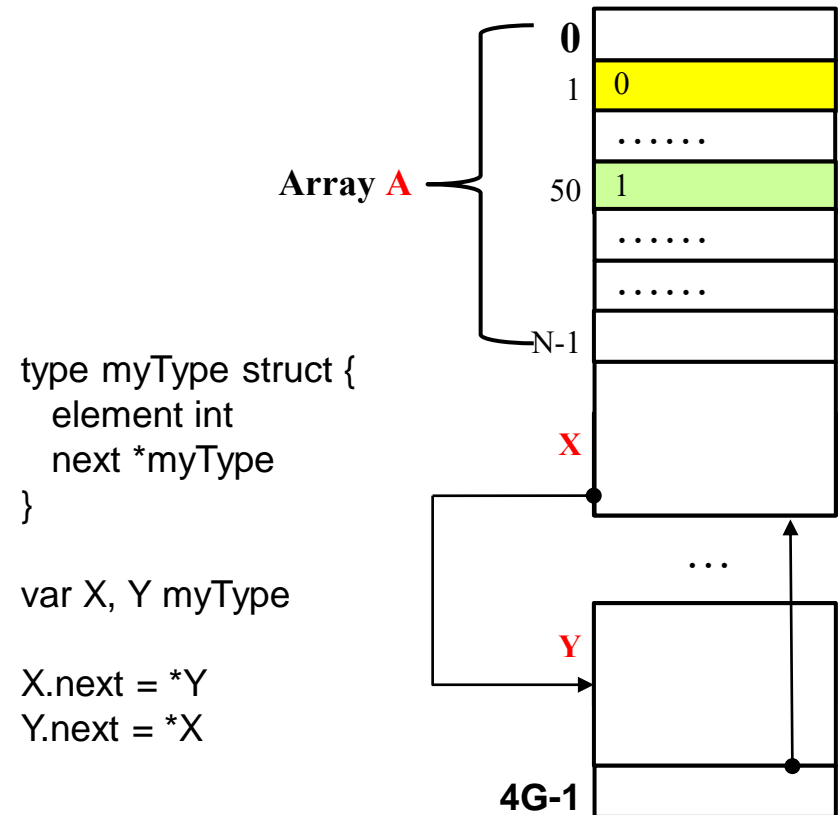
Mask  
mechanism

x = 00111111		Given input
v = 00101010		Given input
v = 00101010 & 00000011	= 000000 <b>10</b>	Bitwise AND
x = 00111111 & 11111100	= 001111 <b>00</b>	Bitwise AND
x = 001111 <b>00</b>   000000 <b>10</b>	= 001111 <b>10</b>	Bitwise OR

Note: 0x3 = 000000**11**; 0xFC = 111111**00**

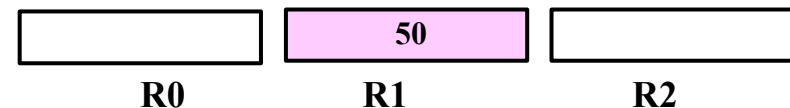
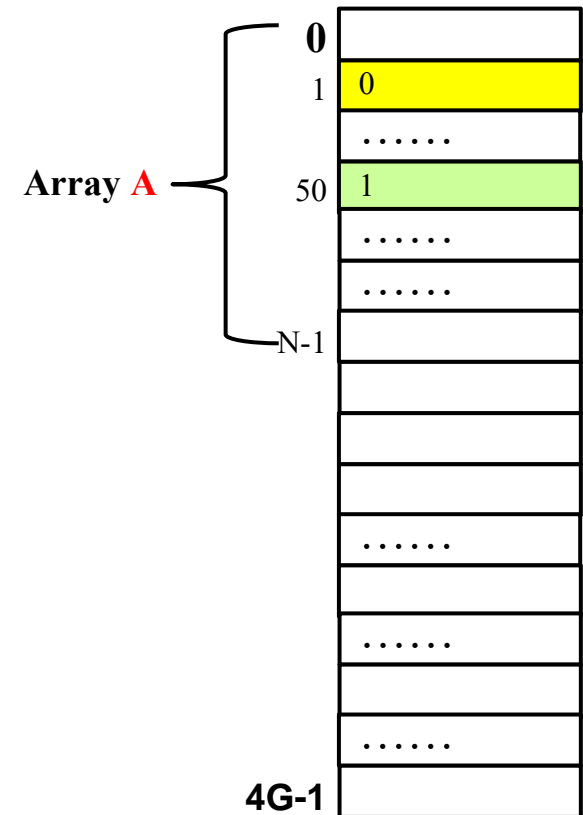
## 3.4.7 Pointers and addressing modes

- Array is simple
  - Linear, consecutive arrangement of N elements of the same type
  - Example: a byte array A
    - Next element of A[i] is A[i+1]
    - If A[i] is at address 50, A[i+1] is at 51
  - What if not consecutive?
- Pointer brings flexibility
  - Nonlinear arrangements, where the elements can jump around
  - Example: two variables X and Y connected by pointers
    - Indicated by the two arrows



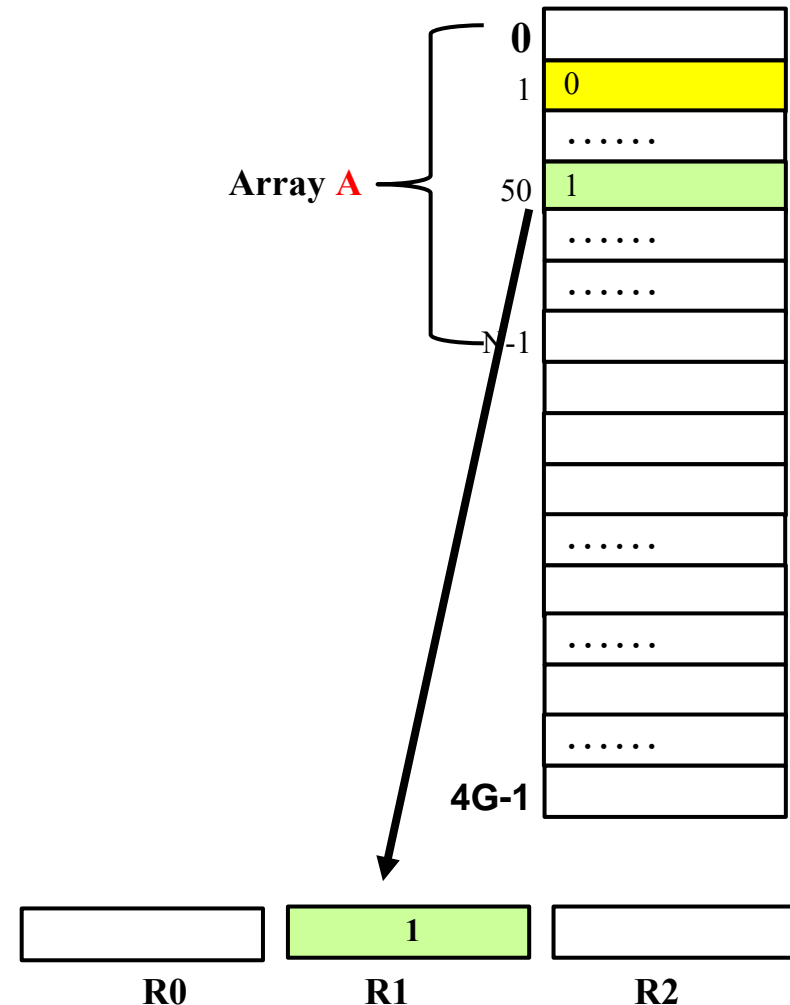
# Contrasting three addressing modes

- Array is simple
  - Linear, consecutive arrangement
- Pointer brings flexibility
  - Nonlinear arrangements
- Pointers are supported by the *indirect addressing mode*
  - Contrasting three addressing modes
    - **Immediate mode:** MOV 50, R1;
      - 50 → R1, i.e., R1=50
    - No memory access



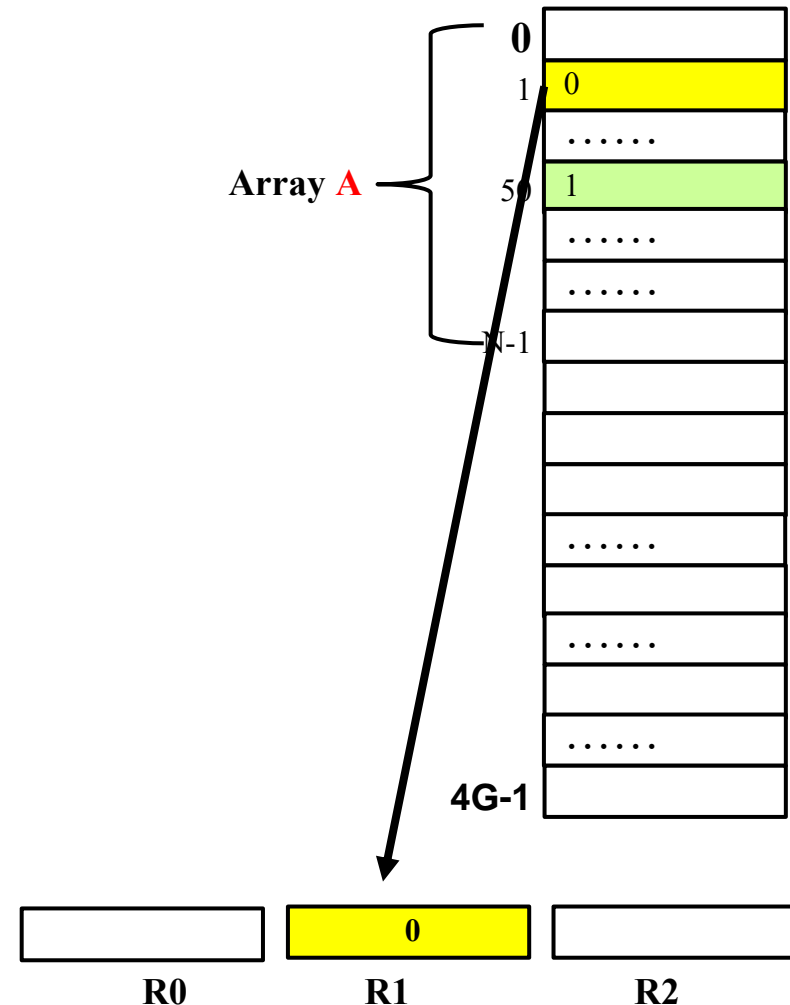
# Contrasting three addressing modes

- Array is simple
  - Linear, consecutive arrangement
- Pointer brings flexibility
  - Nonlinear arrangements
- Pointers are supported by the *indirect addressing mode*
  - Contrasting three addressing modes
    - Immediate mode: MOV 50, R1;
      - 50 → R1, i.e., R1=50
    - **Direct mode**: MOV M[50], R1;
      - M[50] → R1, i.e., R1=1
    - This is the common case



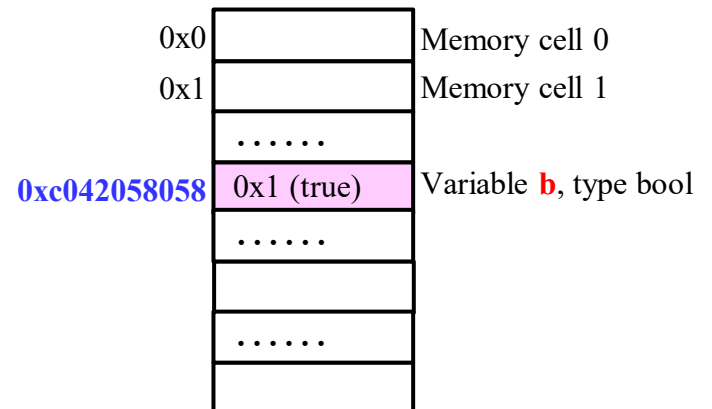
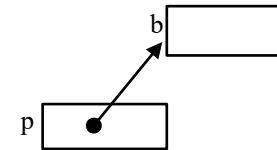
# Contrasting three addressing modes

- Array is simple
  - Linear, consecutive arrangement
- Pointer brings flexibility
  - Nonlinear arrangements
- Pointers are supported by the *indirect addressing mode*
  - Contrasting three addressing modes
    - Immediate mode: `MOV 50, R1;`
      - `50` → R1, i.e., `R1=50`
    - Direct mode: `MOV M[50], R1;`
      - `M[50]` → R1, i.e., `R1=1`
    - Indirect mode: `MOV M[M[50]], R1;`
      - `M[M[50]]` → R1, i.e.,  
`M[M[50]]` is `M[1]`, `R1=M[1]=0`
      - First get the address
      - Then access for the value



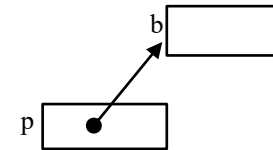
# Step-by-step Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value; dereference p  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```



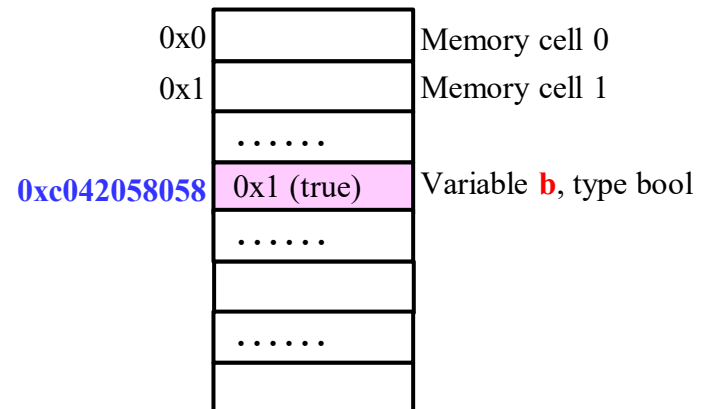
# Step-by-step Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)       // Print b's address  
    fmt.Println(*p)      // Print b's value; dereference p  
    *p = false           // Modify b's value  
    fmt.Println(b)       // Print b's value  
    *p = !(*p)           // Use and modify b's value by negation  
    fmt.Println(b)  
}
```



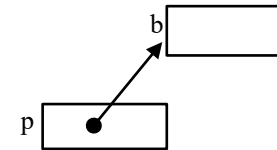
## 1-minute Quiz:

What is the output of the final statement?

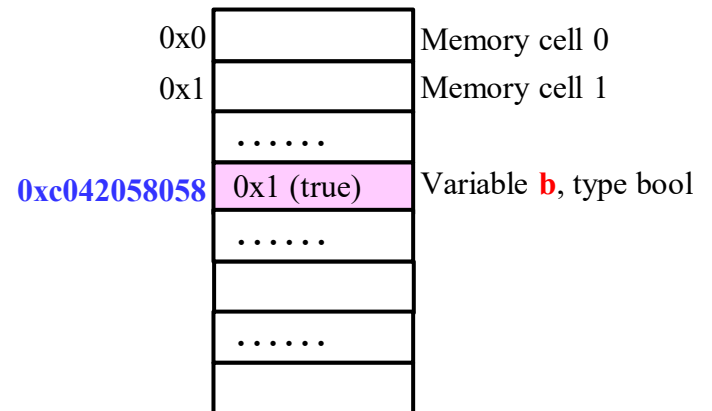


# Step-by-step Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```



```
> go run ./pointer.go  
0xc042058058  
true  
false  
true  
>
```

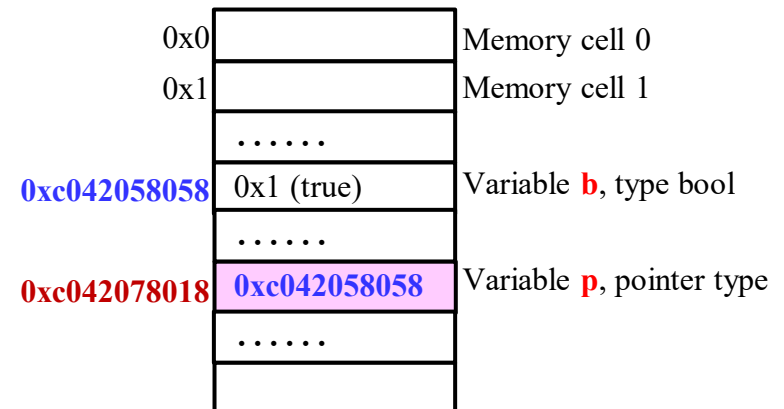
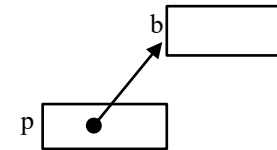




# Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```

```
> go run ./pointer.go  
0xc042058058  
true  
false  
true  
>
```



# Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```

> go run ./pointer.go

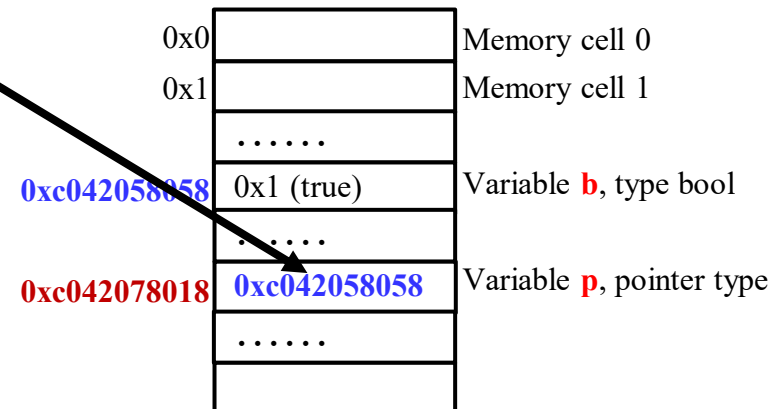
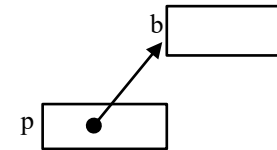
**0xc042058058**

true

false

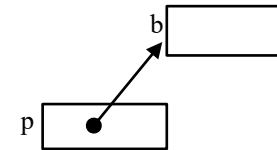
true

>

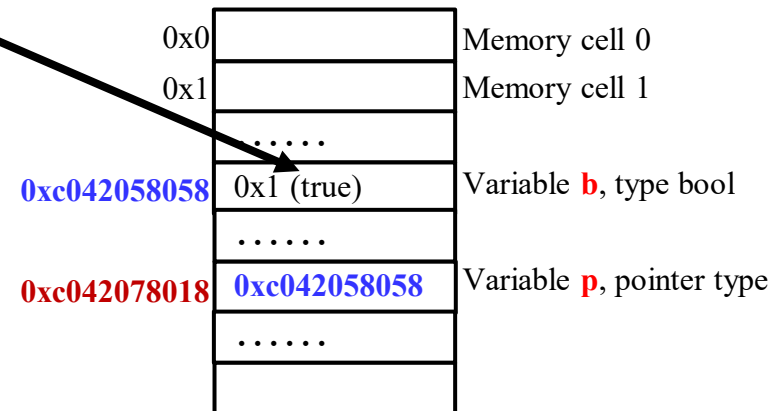


# Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print what *p holds, i.e., b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```



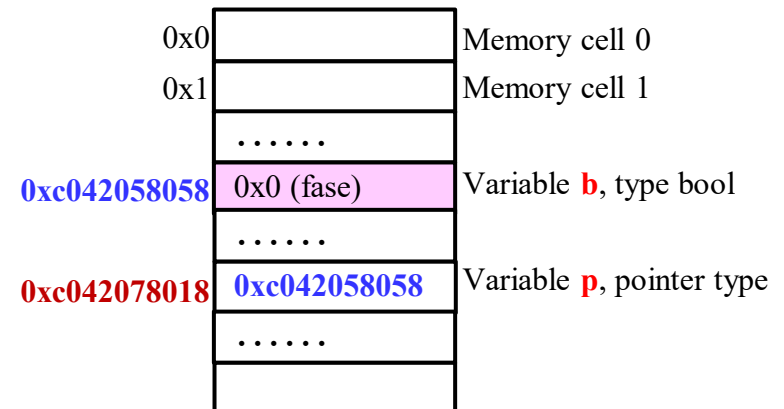
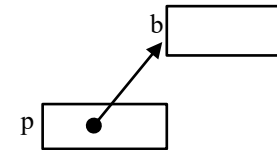
```
> go run ./pointer.go  
0xc042058058  
true  
false  
true  
>
```



# Illustration of pointers

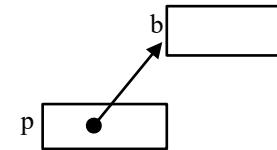
```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```

```
> go run ./pointer.go  
0xc042058058  
true  
false  
true  
>
```

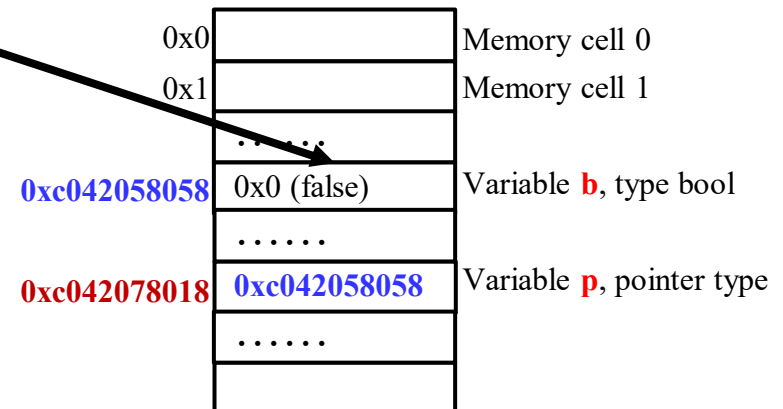


# Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```

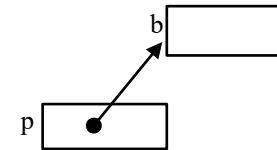


```
> go run ./pointer.go  
0xc042058058  
true  
false  
true  
>
```

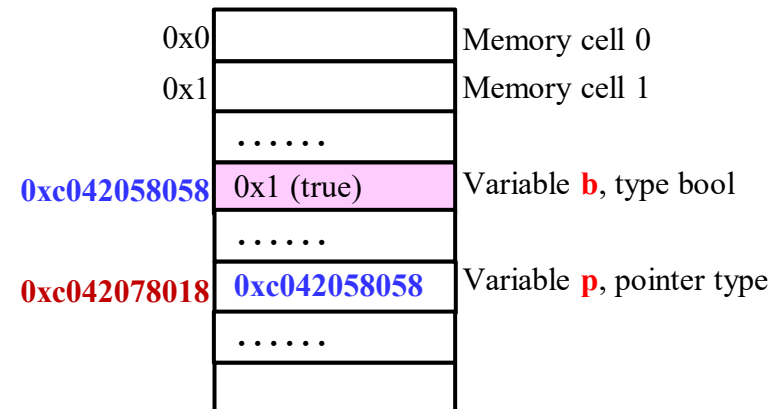


# Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```

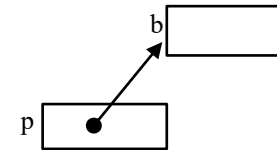


```
> go run ./pointer.go  
0xc042058058  
true  
false  
true  
>
```

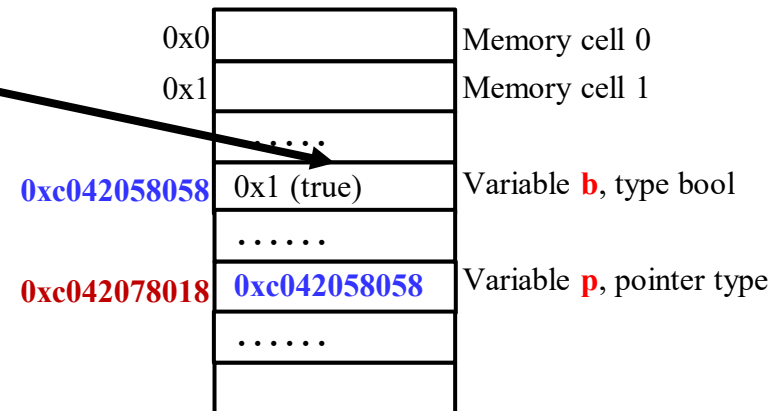


# Illustration of pointers

```
func main() {  
    b := true           // Boolean variable b  
    p := &b             // p holds b's address  
    fmt.Println(p)      // Print b's address  
    fmt.Println(*p)     // Print b's value  
    *p = false          // Modify b's value  
    fmt.Println(b)      // Print b's value  
    *p = !(*p)          // Use and modify b's value by negation  
    fmt.Println(b)  
}
```



```
> go run ./pointer.go  
0xc042058058  
true  
false  
true  
>
```



## 3.4.8 The file abstraction

- To organize and **persistently** store chunks of information
- Files are organized as a hierarchy (tree)
  - Leaf nodes are files; internal nodes are **directories** (special files)
- A file is identified by a **file name** (file path, or **path**)

- **Absolute path**: all the way from the root (/)

- Absolute path of **Autumn.bmp**: /cs101/Prj2/Autumn.bmp

- **Relative path of Autumn.bmp**

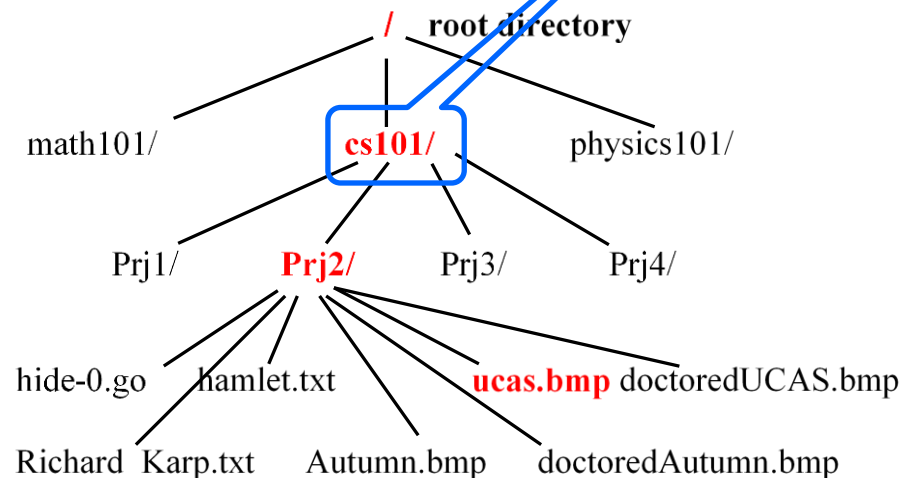
- Related to the **current directory**, i.e., **working directory**

- **./Autumn.bmp** if the working directory is /cs101/Prj2/

- **../Prj2/Autumn.bmp** if the working directory is /cs101/Prj2/

- **Home directory**

- The default directory when login in. Assume /cs101 is the home directory



```
>pwd  
/cs101/
```

You're here



## 3.4.8 The file abstraction

- To organize and **persistently** store chunks of information
- Files are organized as a hierarchy (tree)
  - Leaf nodes are files; internal nodes are **directories** (special files)
- A file is identified by a **file name** (file path, or **path**)

- **Absolute path**: all the way from the root (/)

- Absolute path of Autumn.bmp: /cs101/Prj2/Autumn.bmp

- **Relative path** of Autumn.bmp

- Related to the **current directory**, i.e., **working directory**

- **./Autumn.bmp** if the working directory is /cs101/Prj2/

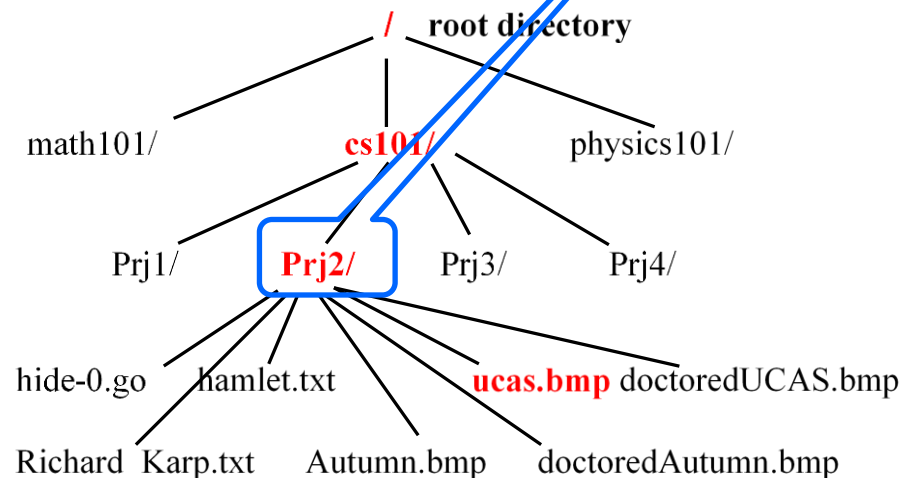
- **../Prj2/Autumn.bmp** if the working directory is /cs101/Prj2/

- **Home directory**

- The default directory when login in. Assume /cs101 is the home directory

```
>cd Prj2
/cs101/Prj2
>pwd
/cs101/Prj2
```

You're here



# Look inside a file: data and metadata

- Data and metadata of file Autumn.bmp
- Data: bits of the actual picture (Pixel Array)
- Metadata: data about the picture data
  - BMP format in the file: File Header, Info Header
  - Other data associated with the file

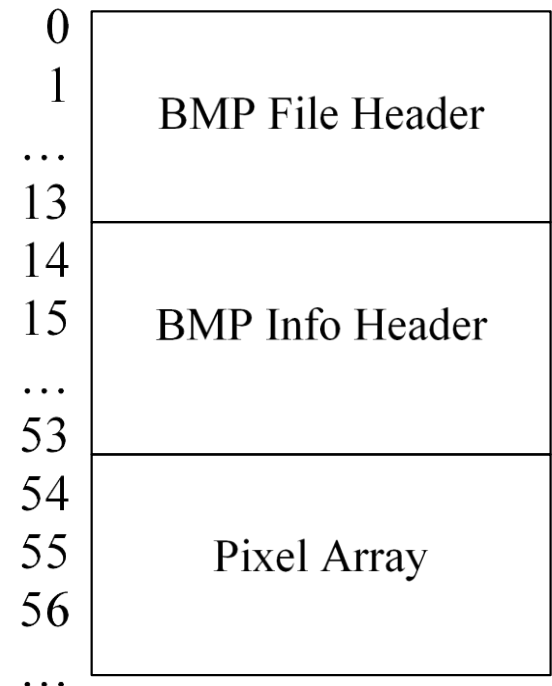
Addresses 0~53 hold metadata

The pixel array for actual image data starts at address 54



Autumn.bmp

The extension **bmp** says it's a bit map image file



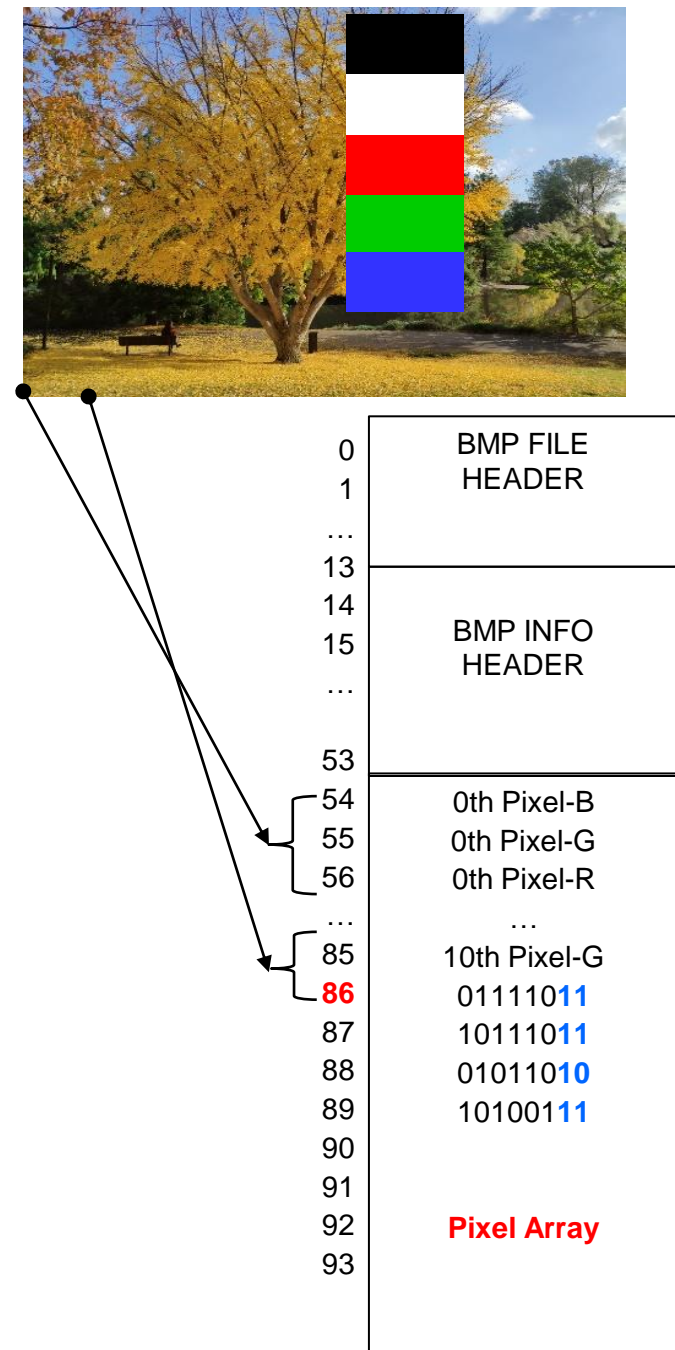
# Other types of metadata

- Can be seen by running “ls -l Autumn.bmp”
  - The file name, the file size, the time of creation (last modification)
  - Access permissions
    - Rights to read, write, and execute a file by the owner of the file, by the group the owner belonging to, and by other users
- Example of access permissions
  - ioutil.WriteFile("./doctoredAutumn.bmp", p, 0666)
  - Every user can to read and write, but cannot execute
    - -rw-rw-rw-
    - 0666 = 0110110110

Owner			Group			Others		
r	w	e	r	w	e	r	w	e
1	1	-	1	1	-	1	1	-

# How is the image of Autumn.bmp stored in Pixel Array?

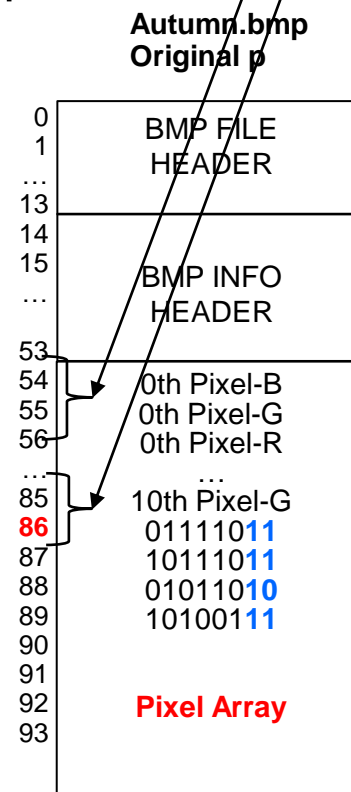
- Pixel = Picture Element
- Pixel Array holds the pixels of the image, starting from the low-left corner going right, and then up row by row
- Each pixel has three bytes for
  - Color depth values of RGB, i.e., the primary colors of red, green, blue
  - RGB values = (0, 0, 0) → Black
  - RGB values = (255, 255, 255) → White
  - RGB values = (255, 0, 0) → Red
  - RGB values = (0, 255, 0) → ?
  - RGB values = (0, 0, 255) → ?
- The first pixel is the 0th element of Pixel Array
  - Uses addresses 54, 55, and 56 to store its three RGB color depth values



# How to hide the length of a text file in a picture?

We need the length when recover in show.go

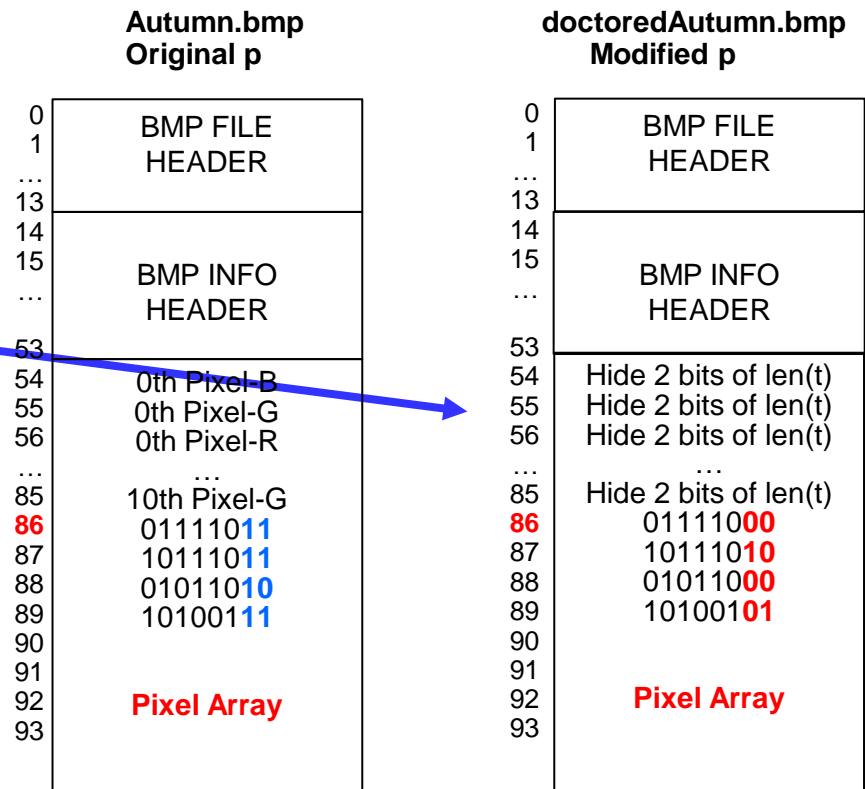
- `p, _ := ioutil.ReadFile("./Autumn.bmp")`  
to read the image file into byte array `p`
- Recall that a function can return multiple values
  - This function returns two values, the second of which is not needed by this code
  - Use a placeholder symbol `'_'`
  - Also called **the blank identifier**





# How to hide the length of a text file in a picture?

- $p, \_ := \text{ioutil.ReadFile}("./\text{Autumn.bmp}")$   
to read the image file into byte array  $p$
- $t, \_ := \text{ioutil.ReadFile}("./\text{hamlet.txt}")$   
to read the text file into byte array  $t$
- Length  $\text{len}(t)$  is a 64-bit integer
  - Hide every 2 bits in a byte of  $p$
  - Need 32 bytes
  - $S = 54, T = 32$
- $\text{modify}(\text{len}(t), p[S:S+T], T)$   
to hide  $\text{len}(t)$  in  $p[54:86]$



# How to hide the length of a text file in a picture?

- `p, _ := ioutil.ReadFile("./Autumn.bmp")`  
to read the image file into byte array `p`
- `t, _ := ioutil.ReadFile("./hamlet.txt")`  
to read the text file into byte array `t`
- Length `len(t)` is a 64-bit integer
  - Hide every 2 bits in a byte of `p`
  - Need 32 bytes
  - $S = 54, T = 32$
- `modify(len(t), p[S:S+T], T)`  
to hide `len(t)` in `p[54:86]`

```
func modify(txt int, pix []byte, size int) {
    for i := 0; i < size; i++ {
        replace last 2 bits of pix[i]
        with the last 2 bits of txt
        repeat with the next 2 bits of txt
    }
}
```

See 3.4.6 in this lecture



Autumn.bmp Original p		doctoredAutumn.bmp Modified p	
0	BMP FILE	0	BMP FILE
1	HEADER	1	HEADER
...		...	
13		13	
14	BMP INFO	14	BMP INFO
15	HEADER	15	HEADER
...		...	
53		53	
54	0th Pixel-B	54	Hide 2 bits of len(t)
55	0th Pixel-G	55	Hide 2 bits of len(t)
56	0th Pixel-R	56	Hide 2 bits of len(t)
...		...	
85	10th Pixel-G	85	Hide 2 bits of len(t)
86	01111011	86	01111000
87	10111011	87	10111010
88	01011010	88	01011000
89	10100111	89	10100101
90		90	
91		91	
92	Pixel Array	92	Pixel Array
93		93	

# How to hide the contents of a text file in a picture?

- $p, \_ := \text{ioutil.ReadFile}("./\text{Autumn.bmp}")$   
to read the image file into byte array  $p$
- $t, \_ := \text{ioutil.ReadFile}("./\text{hamlet.txt}")$   
to read the text file into byte array  $t$
- $t[0]$  holds the **1st character** 'H' = 72
- $\text{modify}(\text{int}(t[0]), p[S+T:S+T+C], C)$   
where
  - $t[0]$  is 'H' = 72 = **01001000**
  - $S = 54$ ,  $T = 32$ ,  $C$  is 4
  - $p[S+T:S+T+C]$  is  $p[86:90]$

Original  $p[86:90]$

86	011110 <b>11</b>
87	101110 <b>11</b>
88	010110 <b>10</b>
89	101001 <b>11</b>



Modified  $p[86:90]$

	011110 <b>00</b>
	101110 <b>10</b>
	010110 <b>00</b>
	101001 <b>01</b>



Autumn.bmp  
Original  $p$

0	BMP FILE HEADER
1	
...	
13	
14	BMP INFO HEADER
15	
...	
53	
54	0th Pixel-B
55	0th Pixel-G
56	0th Pixel-R
...	...
85	10th Pixel-G
<b>86</b>	011110 <b>11</b>
87	101110 <b>11</b>
88	010110 <b>10</b>
89	101001 <b>11</b>
90	
91	
92	<b>Pixel Array</b>
93	

doctoredAutumn.bmp  
Modified  $p$

0	BMP FILE HEADER
1	
...	
13	
14	BMP INFO HEADER
15	
...	
53	
54	Hide 2 bits of len(t)
55	Hide 2 bits of len(t)
56	Hide 2 bits of len(t)
...	...
85	Hide 2 bits of len(t)
<b>86</b>	011110 <b>00</b>
87	101110 <b>10</b>
88	010110 <b>00</b>
89	101001 <b>01</b>
90	
91	
92	<b>Pixel Array</b>
93	



# How to hide the contents of a text file in a picture?

- `p, _ := ioutil.ReadFile("./Autumn.bmp")`  
to read the image file into byte array `p`
- `t, _ := ioutil.ReadFile("./hamlet.txt")`  
to read the text file into byte array `t`
- To hide all `t[i]`,  $i = 0$  to `len(t)`

```
for i:=0; i<len(t); i++{
    offset := S+T+(i*4)
    modify(int(t[i]), p[offset:offset+C], C)
}
```

- Each iteration hides `t[i]` in `p[S+T+(i*4):S+T+(i*4)+C]`
  - Where  $S = 54$ ,  $T = 32$ ,  $C = 4$
- That is, `t[i]` is hidden in `p[86+(i*4)]`, `p[86+(i*4)+1]`, `p[86+(i*4)+2]`, `p[86+(i*4)+3]`
- E.g., `t[1]='A'`, is hidden in `p[90:94]`



	Autumn.bmp Original p	doctoredAutumn.bmp Modified p
0	BMP FILE	BMP FILE
1	HEADER	HEADER
...		
13		
14	BMP INFO	BMP INFO
15	HEADER	HEADER
...		
53		
54	0th Pixel-B	Hide 2 bits of len(t)
55	0th Pixel-G	Hide 2 bits of len(t)
56	0th Pixel-R	Hide 2 bits of len(t)
...		
85	10th Pixel-G	Hide 2 bits of len(t)
86	01111011	01111000
87	10111011	10111010
88	01011010	01011000
89	10100111	10100101
90		
91		
92	Pixel Array	Pixel Array
93		

# Check results

- Write the complete `hide-0.go`
- Execute and display result
  - > `go run hide-0.go`
  - > `display doctoredAutumn.bmp`
- The Text Hider project
  - Produce `hide.go` with good coding practices
  - Also need to write `show.go`
- Change `hide-0.go` to `hide-1.go`
  - by modifying the most significant 2 bits (the rightmost 2 bits) of each byte of Pixel Array
  - > `go run hide-1.go`
  - > `display doctoredAutumn.bmp`



Original Autumn.bmp



doctoredAutumn.bmp  
Modifying rightmost 2 bits



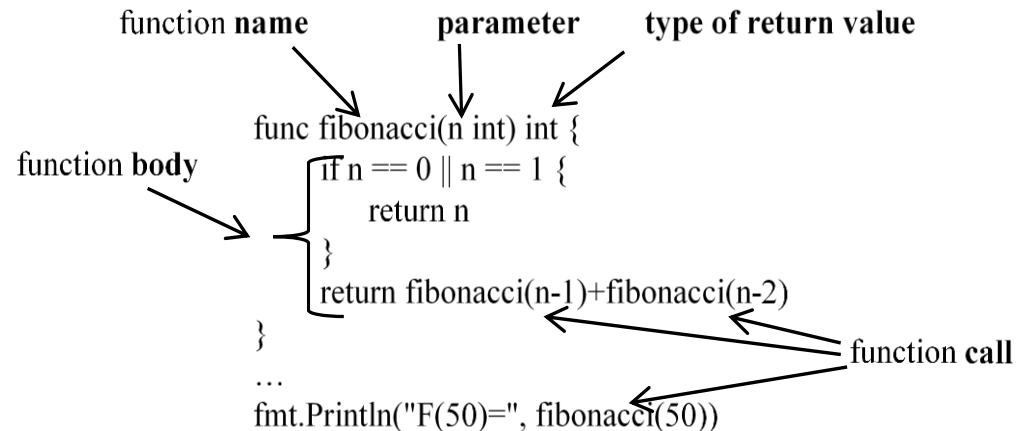
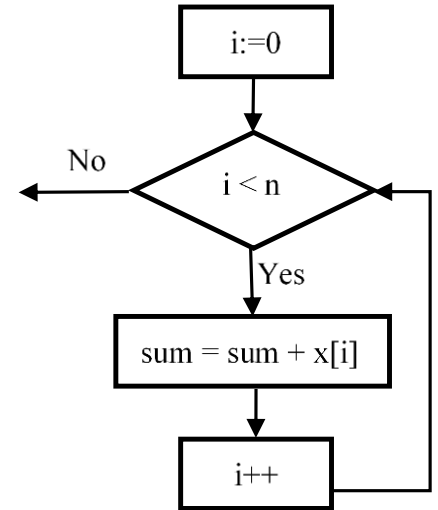
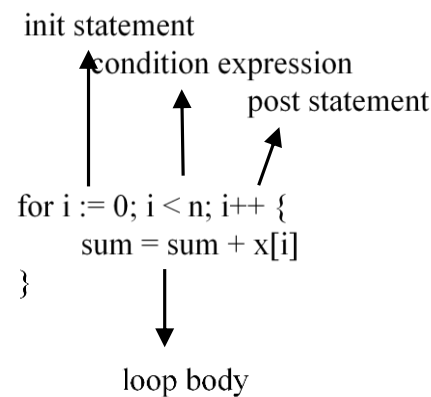
doctoredAutumn.bmp  
Modifying leftmost 2 bits

# 3.5 Review of control abstractions

- **Precedence** in an expression
  - For  $x*b+c \parallel i < 7$ , the precedence is  $((x*b)+c) \parallel (i < 7)$
  - When in doubt, use parentheses
- **Sequence** of statements: follow the syntactic order
- **Selection** (conditional): if-then-else statement

```
if i<7 {  
    fmt.Println(i)  
}
```

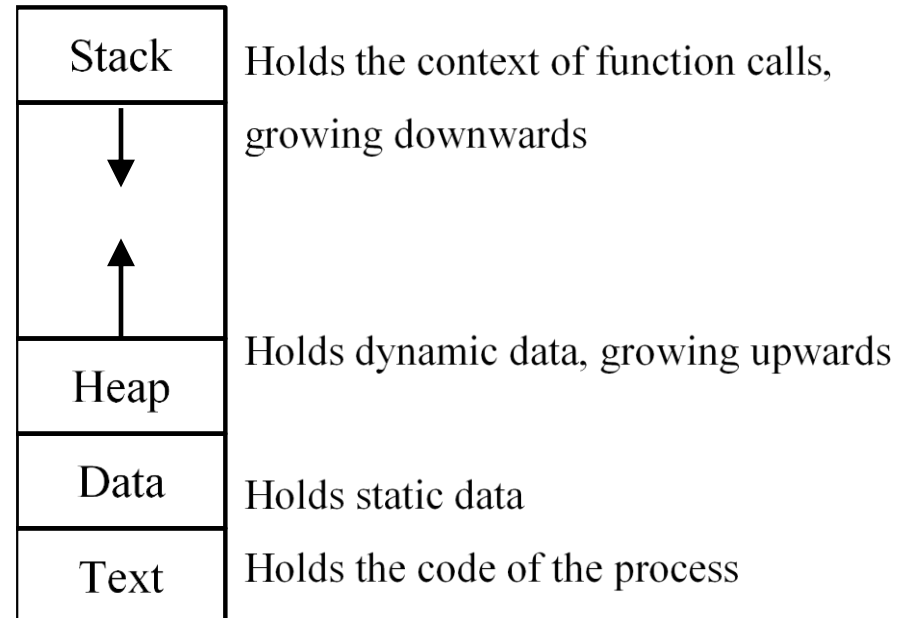
- **Loop**: repetitively execute a body of code
- **Function**: defined once and can be called many times



# Segments of Text, Data, Heap and Stack

- Segments: areas of memory
  - Text: code of fib-50.go
  - Data: static data
  - Heap:
  - Stack: context of function calls
- How much memory is needed for computing  $F(5000000)$ ?
  - A small constant? Multiple of 5 million?  $2^{5000000}$ ?

```
package main
import "fmt"
func main() {
    fmt.Println("F(50)=", fibonacci(50))
}
func fibonacci(n int) int {
    if n == 0 || n == 1 {
        return n
    }
    return fibonacci(n-1)+fibonacci(n-2)
}
```



# \*\*\* Computing Fibonacci numbers of arbitrary word length, with fib.Uint.go

```
package main // fib.Uint.go
import (
    "fmt"
    "math"
)
func main() {
    fmt.Printf("F(100) = %s\n", String(*(fibonacci(100))))
}
type Uint []uint64
func fibonacci(n int) *Uint {
    a := &Uint{0} // a = 0
    b := &Uint{1} // b = 1
    for i := 1; i < n+1; i++ {
        Acc(a, b) //a = a + b
        a, b = b, a
    }
    return a
}
// Code for func Acc() and func String() functions
```

```
> go run fib.Uint.go
F(100) = 354224848179261915075
>
```

Program fib.Uint.go uses a slice of uint64 numbers to represent an unsigned integer of arbitrary length, by defining a type Uint

Function A(a, b) does an accumulation  $a = a + b$  where a and b are of type \*Uint

```
func Acc(a, b *Uint) {
    x, y := *a, *b
    d := make(Uint, len(y)-len(x))
    x = append(x, d...)
    c := make(Uint, len(x)+1)
    for i := 0; i < len(x); i++ {
        var v uint64
        v = x[i] + y[i] + c[i]
        if v < x[i] || v < y[i] || v < c[i] {
            c[i+1] = 1
        }
        c[i] = v
    }
    if c[len(c)-1] == 0 {
        c = c[:len(c)-1]
    }
    *a = c
}
```

Students donot need to  
Know the internal of Acc

## \*\*\* Code of fib.Uint.go, continued

```
package main // fib.Uint.go
import (
    "fmt"
    "math"
)
func main() {
    fmt.Printf("F(100) = %s\n", String(*(fibonacci(100))))
}
type Uint []uint64
func fibonacci(n int) *Uint {
    a := &Uint{0}           // a = 0
    b := &Uint{1}           // b = 1
    for i := 1; i < n+1; i++ {
        Acc(a, b)           //a = a + b
        a, b = b, a
    }
    return a
}
// Code for func Acc() and func String() functions
```

```
> go run fib.Uint.go
F(100) = 354224848179261915075
>
```

String(x) converts a Uint value into a decimal string

```
func String(x Uint) string {
    if len(x) == 1 && (x)[0] == 0 {
        return "0"
    }
    i := int(float64(len(x)*64)/math.Log2(float64(10))) + 1
    s := make([]byte, i)
    var r byte = 0
    for ; len(x) != 0; i-- {
        temp := make(Uint, len(x))
        var dividend uint64 = 0
        for i := len(x) - 1; i >= 0; i-- {
            dividend = dividend<<32 + uint64(x[i]>>32)
            q := dividend / 10
            r = byte(dividend - (q<<3 + q<<1))
            temp[i] = q
            dividend = uint64(r)
            dividend = dividend<<32 + uint64(x[i]<<32>>32)
            q = dividend / 10
            r = byte(dividend - (q<<3 + q<<1))
            temp[i] = temp[i]<<32 + q
            dividend = uint64(r)
        }
        if temp[len(temp)-1] == 0 {
            temp = temp[:len(temp)-1]
        }
        x, s[i-1] = temp, r+'0'
    }
    i = 0
    for s[i] == 0 { i++ }
    return string(s[i:])
}
```

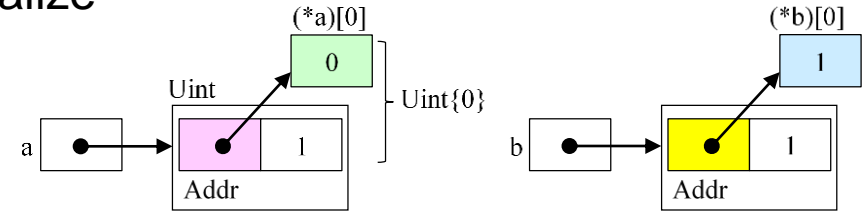
Students donot need to  
Know the internal of String

# \*\*\* Some execution details

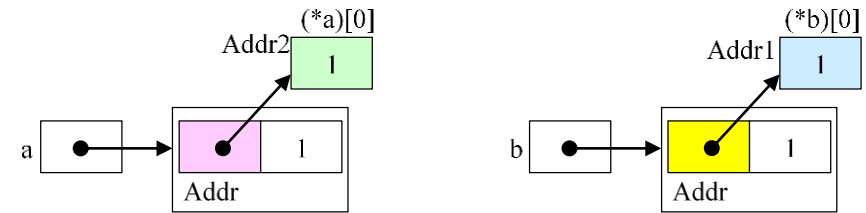
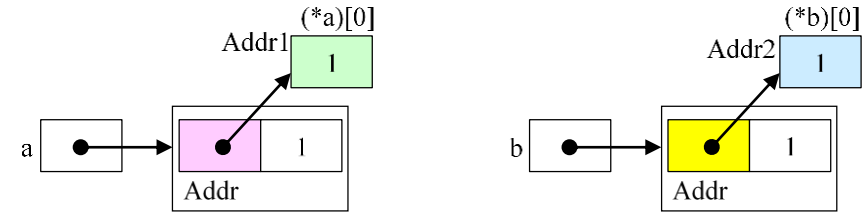
```
package main // fib.Uint.go
import (
    "fmt"
    "math"
)
func main() {
    fmt.Printf("F(100) = %s\n", String(*(fibonacci(100))))
}
type Uint [uint64]
func fibonacci(n int) *Uint {
    a := &Uint{0} // a = 0
    b := &Uint{1} // b = 1
    for i := 1; i < n+1; i++ {
        Acc(a, b) // a = a + b
        a, b = b, a
    }
    return a
}
// Code for func Acc() and func String() functions
```

```
> go run fib.Uint.go
F(100) = 354224848179261915075
>
```

Initialize



When i = 1

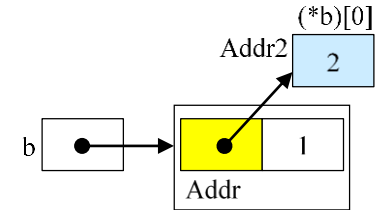
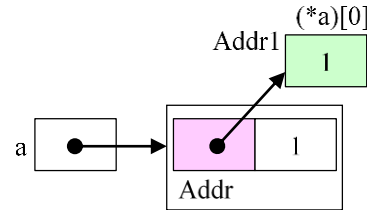
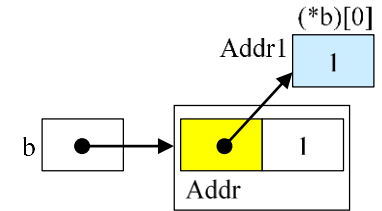
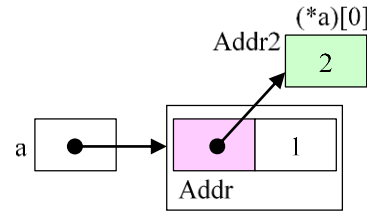




# \*\*\* Some execution details

```
package main // fib.Uint.go
import (
    "fmt"
    "math"
)
func main() {
    fmt.Printf("F(100) = %s\n", String(*(fibonacci(100))))
}
type Uint [uint64]
func fibonacci(n int) *Uint {
    a := &Uint{0} // a = 0
    b := &Uint{1} // b = 1
    for i := 1; i < n+1; i++ {
        Acc(a, b) // a = a + b
        a, b = b, a
    }
    return a
}
// Code for func Acc() and func String() functions
```

When i = 2



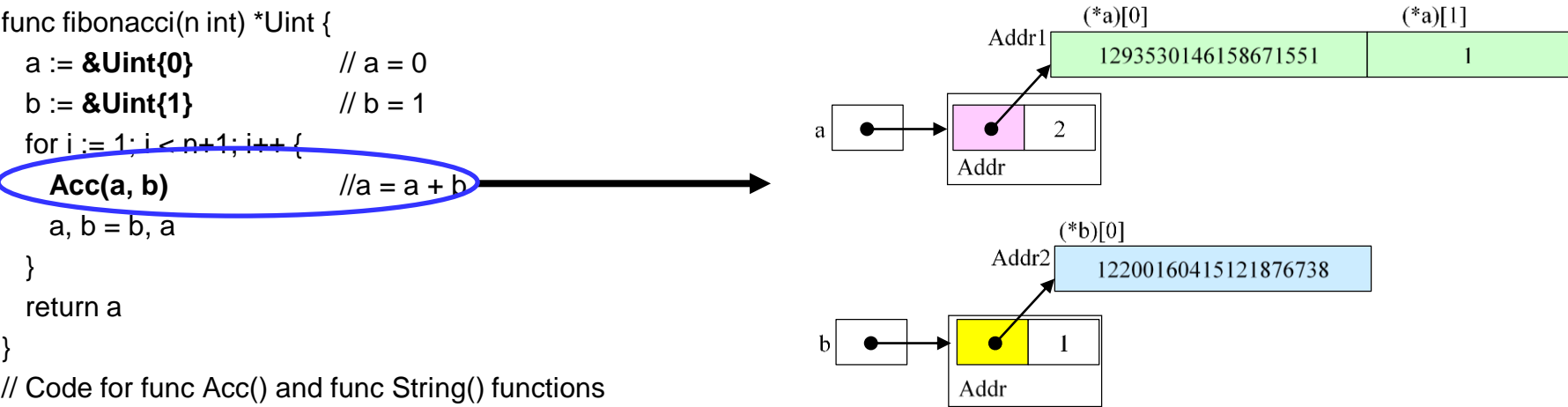
```
> go run fib.Uint.go
F(100) = 354224848179261915075
>
```



# \*\*\* Some execution details

```
package main // fib.Uint.go
import (
    "fmt"
    "math"
)
func main() {
    fmt.Printf("F(100) = %s\n", String(*(fibonacci(100))))
}
type Uint []uint64
func fibonacci(n int) *Uint {
    a := &Uint{0} // a = 0
    b := &Uint{1} // b = 1
    for i := 1; i < n+1; i++ {
        Acc(a, b) //a = a + b
        a, b = b, a
    }
    return a
}
// Code for func Acc() and func String() functions
```

When i = 93



```
> go run fib.Uint.go
F(100) = 354224848179261915075
>
```